

Figure 13.1 A recursive 8-point unordered FFT computation.



**Figure 13.2** The pattern of combination of elements of the input and the intermediate sequences during a 16-point unordered FFT computation.



**Figure 13.3** A 16-point unordered FFT on 16 processes. P<sub>i</sub> denotes the process labeled *i*.



**Figure 13.4** A 16-point FFT on four processes.  $P_i$  denotes the process labeled *i*. In general, the number of processes is  $p = 2^d$  and the length of the input sequence is  $n = 2^r$ .



**Figure 13.5** Isoefficiency functions of the binary-exchange algorithm on a hypercube with  $t_c = 2$ ,  $t_w = 4$ , and  $t_s = 25$  for various values of *E*.



**Figure 13.6** The efficiency of the binary-exchange algorithm as a function of *n* on a 256-node hypercube with  $t_c = 2$ ,  $t_w = 4$ , and  $t_s = 25$ .



**Figure 13.7** Data communication during an FFT computation on a logical square mesh of 64 processes. The figure shows all the processes with which the processes labeled 0 and 37 exchange data.



**Figure 13.8** The pattern of combination of elements in a 16-point FFT when the data are arranged in a  $4 \times 4$  two-dimensional square array.



(a) Steps in phase 1 of the transpose algorithm (before transpose)



(b) Steps in phase 3 of the transpose algorithm (after transpose)

Figure 13.9 The two-dimensional transpose algorithm for a 16-point FFT on four processes.



**Figure 13.10** Data distribution in the three-dimensional transpose algorithm for an *n*-point FFT on p processes ( $\sqrt{p} \le n^{1/3}$ ).



**Figure 13.11** The communication (transposition) phases in the three-dimensional transpose algorithm for an n-point FFT on p processes.



**Figure 13.12** A comparison of the speedups obtained by the binary-exchange, 2-D transpose, and 3-D transpose algorithms on a 64-node hypercube with  $t_c = 2$ ,  $t_w = 4$ , and  $t_s = 25$ .