

Data Structures for Sets

- Many applications deal with sets.
 - Compilers have symbol tables (set of vars, classes)
 - Dictionary is a set of words.
 - Routers have sets of forwarding rules.
 - Web servers have set of clients, etc.
- A set is a collection of members
 - No repetition of members
 - Members themselves can be sets
- Examples
 - Set of first 5 natural numbers: $\{1,2,3,4,5\}$
 - $\{x \mid x \text{ is a positive integer and } x < 100\}$
 - $\{x \mid x \text{ is a CA driver with } > 10 \text{ years of driving experience and } 0 \text{ accidents in the last 3 years}\}$

Set Operations

Binary operations	Member	Set
Member	Order (=, <, >)	Find, insert, delete, split, ...
Set	Find, insert, delete, split, ...	Union, intersection, difference, equal, ...

- Unary operation: min, max, sort, makenull, ...

Observations

- Set + Operations define an ADT.
 - A set + insert, delete, find
 - A set + ordering
 - Multiple sets + union, insert, delete
 - Multiple sets + merge
 - Etc.
- Depending on type of members and choice of operations, different implementations can have different asymptotic complexity.

Set ADT: Union, Intersection, Difference

AbstractDataType *SetUID*

instance
multiple sets

operations

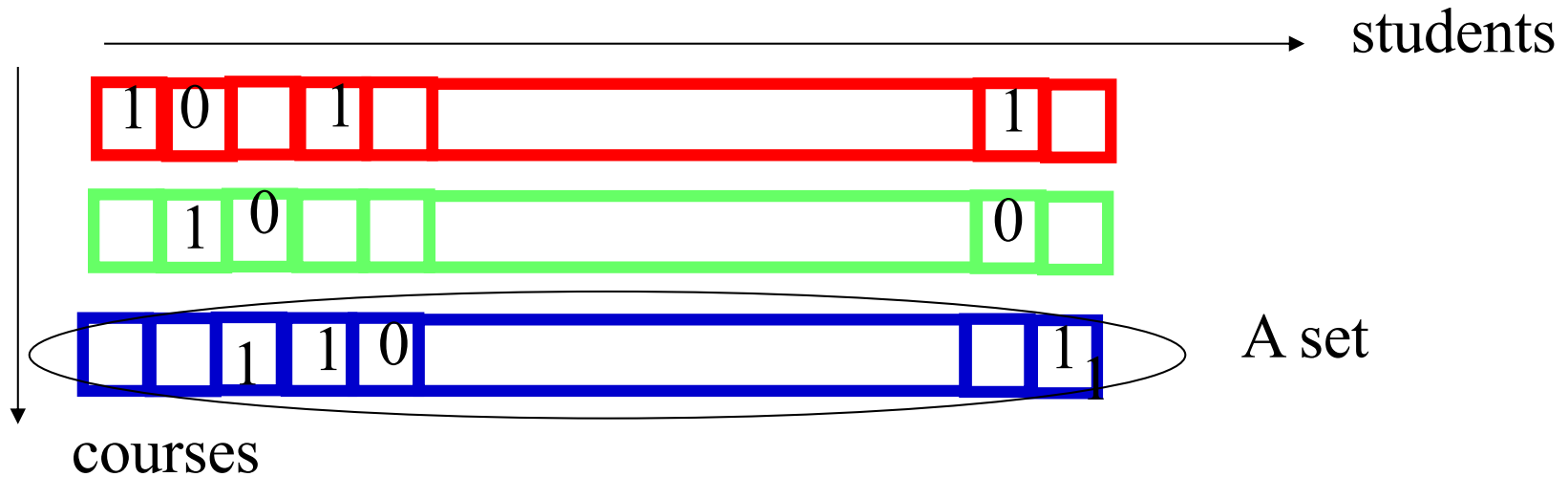
union (*s1*, *s2*): $\{x \mid x \text{ in } s1 \text{ or } x \text{ in } s2\}$
intersection (*s1*, *s2*): $\{x \mid x \text{ in } s1 \text{ and } x \text{ in } s2\}$
difference (*s1*, *s2*): $\{x \mid x \text{ in } s1 \text{ and } x$
not in s2\}

Examples

- **Sets:** Articles in Yahoo Science (A), Technology (B), and Sports (C)
 - Find all articles on Wright brothers.
 - Find all articles dealing with sports medicine
- **Sets:** Students in CS8 (A), CS16 (B), and CS40 (C)
 - Find all students enrolled in these courses
 - Find students registered for CS8 only
 - Find students registered for both CS8 and CS16
 - Etc.

Set UID Implementation: Bit Vector

- Set members known and finite (e.g., all students in CS dept)



■ Operations

- Union: $u[k] = x[k] \mid y[k]$;
- Intersection: $u[k] = x[k] \& y[k]$;
- Difference: $u[k] = x[k] \& \sim y[k]$;

- Complexity: $O(n)$: n size of the set

Set UID Implementation: linked lists

■ Bit vectors great when

- Small sets
- Known membership

■ Linked lists

- Unknown size and members
- Two kinds: Sorted and Unsorted

Set UID Complexity: **Unsorted** Linked List

■ Intersection

For $k=1$ to n do

Advance set_A one step to find k th element;

Follow set_B to find that element in B ;

If found then

Append element k to set_{AB}

End

■ Searching for each element can take n steps.

■ Intersection worst-case time $O(n^2)$.

Set UID Complexity: **Sorted** Lists

- The list is sorted; larger elements are to the right
- Each list needs to be scanned only once.
- At each element: increment and possibly insert into A&B, constant time operation
- Hence, sorted list set-set ADT has $O(n)$ complexity
- A simple example of how even trivial algorithms can make a big difference in runtime complexity.

Set UID: Sorted List Intersection

- **Case A** $*set_A = *set_B$
 - Include $*set_A$ (or $*set_B$) in $*set_{AB}$
 - Increment set_A
 - Increment set_B
- **Case B** $*set_A < *set_B$
 - Increment set_A Until
 - $*set_A = *set_B$ (A)
 - $*set_A > *set_B$ (C)
 - $*set_A == null$
- **Case C** $*set_A > *set_B$
 - Increment set_B Until
 - $*set_A = *set_B$ (A)
 - $*set_A < *set_B$ (B)
 - $*set_B == null$
- **Case D** $*set_A == null$ or $*set_B == null$
 - terminate

Dictionary ADTs

- Maintain a set of items with distinct keys with:
 - *find*(k): find item with key k
 - *insert*(x): insert item x into the dictionary
 - *remove*(k): delete item with key k
- Where do we use them:
 - Symbol tables for compiler
 - Customer records (access by name)
 - Games (**positions, configurations**)
 - Spell checkers
 - Peer to Peer systems (access songs by name), etc.

Naive Implementations

- The simplest possible scheme to implement a dictionary is “log file” or “audit trail”.
 - Maintain the elements in a linked list, with insertions occurring at the head.
 - The search and delete operations require searching the entire list in the worst-case.
 - Insertion is $O(1)$, but find and delete are $O(n)$.
- A sorted array does not help, even with ordered keys. The search becomes fast, but **insert/delete take $O(n)$** .